

# Assignment 2

What you are allowed to do:

- Discuss the problems with your classmates.
- Ask the teacher for help (some points are hard, it is ok not to know how to proceed).
- Use a computer for linear operations (e.g. matrix multiplication, matrix vector etc), provided you know how to do it manually.

What you are bound to do:

- Write down the solutions yourself, while being alone.

What you are not allowed to do:

- Copy the written solutions from a classmate.
- Use AI to output the solutions.

## Problems

**Exercise 0.1.** Let  $P$  be the Markov operator of an irreducible Markov Chain on a finite state space  $S$ . Let  $m$  be its invariant measure, and  $P^\dagger$  the adjoint operator of  $P$  in  $L^2(m)$ .

- Check that if  $P, Q$  are Markov operators, then  $PQ$  and  $\alpha P + (1 - \alpha)Q$  are Markov operators. In particular  $PP^\dagger$  and  $P^2$  are Markov operators.
- Find the invariant measure of  $PP^\dagger$  and  $P^2$ .
- Give an example where  $P$  is irreducible, but  $P^2$  and  $PP^\dagger$  are not irreducible.
- Prove that, if  $P$  is irreducible and aperiodic, then  $P^2$  is irreducible and aperiodic.
- Prove that  $PP^\dagger$  is aperiodic and each communicating class is closed.
- Assume that  $P$  is irreducible and aperiodic. Let  $X_t$  and  $Y_t$  be Markov Chains with Markov operators  $P^2$  and  $PP^\dagger$  respectively, prove that  $X_t$  converges to its “equilibrium” faster than  $Y_t$  in the following sense. Let  $\mu_t^x$  (respectively  $\nu_t^x$ ) be the law of  $X_t$  when  $X_0 = x$  (respectively  $Y_t$  when  $Y_0 = x$ ). Prove that

$$\sup_x \lim_{t \rightarrow \infty} \frac{1}{t} \log \|\mu_t^x - m\|_{TV} \leq \sup_x \lim_{t \rightarrow \infty} \frac{1}{t} \log \|\nu_t^x - m\|_{TV}$$

## Solution

- Let  $R = PQ$ , thus  $r_{x,z} = \sum_y p_{x,y} q_{y,z}$ . And exchanging summations  $r_{x,z} \geq 0$  and  $\sum_z r_{x,z} = \sum_y p_{x,y} \sum_z q_{y,z} = \sum_y p_{x,y} = 1$ . As for the convex combination, the proof is immediate by linearity.
- The invariant measure of all these operators is  $m$ . In general, if  $P, Q$  have the same invariant measure, then  $m((PQ)f) = m(P(Qf)) = m(Qf) = m(f)$ . So  $m$  is invariant for  $PQ$ .
- We can take a unique cycle  $S = \mathbb{Z}_{2k}$  with an even number of points,  $P$  corresponds to moving clockwise with probability 1,  $p_{x,x+1} = 1$  (sums are understood  $(\text{mod } 2)k$ ). Then  $P$  is irreducible since there is a

path joining any two points.  $PP^\dagger$  corresponds to making one step clockwise and one counter-clockwise, that is  $PP^\dagger = I$ , the walk just does not move.  $P^2$  corresponds to making two steps clockwise  $p_{x,x+2}^{(2)} = 1$ . So there is not path joining points with different parity.

- d. For an irreducible aperiodic chain on a finite state space, there exists  $T$  large enough so that  $p_{x,y}^{(t)} > 0$  for all  $t > T$  and  $x, y \in S$ . In particular the same holds true, changing  $T$  to  $T' = \lceil T/2 \rceil$ , for the entries of  $P^2$ .
- e. Let  $(q_{x,y})$  denote the entries of  $Q = PP^\dagger$ .  $q_{x,z} > 0$  iff there exists  $y \in S$  such that  $p_{x,y}, p_{z,y} > 0$ . In particular  $q_{x,x} > 0$  iff  $q_{z,x} > 0$  and  $q_{x,x} > 0$  for all  $x \in S$  (since each point  $x$  has at least one successor  $y_x$  such that  $p_{x,y_x} > 0$ ). Therefore each  $Q$ -communicating class is closed.
- f. By the [exponential convergence theorem](#), we need to check that  $|\lambda_1(P^2)| \leq |\lambda_1(PP^\dagger)|$ , where  $\lambda_1(Q)$  is the largest (by modulus) eigenvalue of an irreducible Markov operator  $Q$ , except  $\lambda_0 = 1$ ; or equivalently, the largest eigenvalue of  $Q$  restricted to functions  $f$  with  $m(f) = 0$ , where  $m$  is the invariant measure of  $Q$ .

For  $P^2$  we have  $\lambda_1(P^2) = \lambda_1(P)^2 = \overline{\lambda_1(P^\dagger)}^2$ . Take  $f_1$  an associated (complex) eigenfunction (with  $m(f_1) = 0$ ) of  $P^\dagger$ , so that  $P^\dagger f_1 = \lambda_1(P^\dagger) f_1$ . Then, since  $PP^\dagger$  is symmetric,  $\lambda_1(PP^\dagger) > 0$  and thus denoting  $\langle \cdot, \cdot \rangle$  the Hermitian product in  $L^2(m)$

$$|\lambda_1(PP^\dagger)| = \sup_{f: m(f)=0} \frac{\langle f, PP^\dagger f \rangle}{m(|f|^2)} = \sup_{f: m(f)=0} \frac{\langle P^\dagger f, P^\dagger f \rangle}{m(|f|^2)} \geq \frac{\langle P^\dagger f_1, P^\dagger f_1 \rangle}{m(|f_1|^2)} = \overline{\lambda_1(P^\dagger)} \lambda_1(P^\dagger) = |\lambda_1(P^2)|$$

**Exercise 0.2.** Consider an irreducible Markov Chain on a finite state space  $S$  with transition probabilities  $p_{x,y}^0$ . Let  $c(x, y) \in \mathbb{R}$  be defined for all edges  $(x, y)$  such that  $p_{x,y} > 0$ . Fix  $\beta > 0$  and define the tilted transition probabilities

$$p_{x,y}^\beta := \frac{e^{-\beta c(x,y)}}{Z_x^\beta} p_{x,y}^0$$

- a. Determine the value of  $Z_x^\beta$  and check that the Markov Chain with transition probabilities  $p_{x,y}^\beta$  is still irreducible.
- b. Explain why we can assume that  $\min_{y: p_{x,y}^0 > 0} c(x, y) = 0$  and verify that in this case

$$\inf_y p_{x,y}^0 \leq Z_x^\beta \leq 1$$

- c. Let  $m^\beta$  be the invariant probability measure associated to the transition probabilities  $p^\beta$ . Characterize the limit  $\lim_{\beta \rightarrow \infty} m^\beta$ . *Hint: For the sake of simplicity, assume that the values  $\sum_{e \in \theta} c(e)$  are distinct for distinct  $\theta \in \Theta$ , where  $\Theta$  is the space of [rooted spanning in-arborescences](#).*
- d. Give an explicit answer in the case  $c(x, y) = (V(y) - V(x))^+$ , where  $V: S \rightarrow \mathbb{R}$  is a given injective function, and for  $a \in \mathbb{R}$ ,  $a^+ = \max(a, 0)$ . *Hint: For the sake of simplicity, assume that  $p_{x,y}^0 > 0$  whenever  $p_{y,x}^0 > 0$ .*

### Solution

- a. Since  $p_{x,\cdot}^\beta$  is a probability, we have

$$Z_x^\beta := \sum_y e^{-\beta c(x,y)} p_{x,y}^0$$

The Markov chain is still irreducible since  $p_{x,y}^\beta > 0$  iff  $p_{x,y}^0 > 0$ .

- b. If we change  $c(x, y)$  to  $c(x, y) - a(x)$ , then  $p_{x,y}^\beta$  does not change, regardless of  $a(x)$ . Therefore, to

make the computation simpler, we can take  $a(x) = \min_y c(x, y)$ . In other words, we just assume  $\min_y c(x, y) = 0$ . In particular

$$\inf_y p_{x,y}^0 \leq \max_y e^{-\beta c(x,y)} p_{x,y}^0 \leq Z_x^\beta \leq \sum_y e^{-0} p_{x,y}^0 = 1$$

c. We know that

$$m_x^\beta = c^\beta \sum_{\theta \in \Theta_x} w^\beta(\theta), \quad w^\beta(\theta) := \prod_{e \in \theta} p_e^\beta$$

where  $c^\beta > 0$  is the suitable normalization constant (which is just the same sum over  $\theta \in \Theta$ ). Using the definition of  $p_{x,y}^\beta$

$$m_x^\beta = c^\beta \sum_{\theta \in \Theta_x} \frac{e^{-\beta \sum_{e \in \theta} c(e)}}{\prod_{y \neq x} Z_y^\beta} w^0(\theta)$$

where the product in the denominator runs over all the  $y \in S$  from which there is an outgoing edge, namely all the  $y \in S$  except the root  $x$ . For  $Z^\beta := \prod_x Z_x^\beta$  and  $c(\theta) := \sum_{e \in \theta} c(e)$ :

$$m_x^\beta = \frac{c^\beta}{Z^\beta} \sum_{\theta \in \Theta_x} Z_x^\beta e^{-\beta c(\theta)} w^0(\theta)$$

We know that for an irreducible chain  $\Theta_x$  is non-empty for any  $x \in S$ . And simply estimating the sum with the largest summand (as above)

$$Z_x^\beta \max_{\theta \in \Theta_x} e^{-\beta c(\theta)} w^0(\theta) \leq m_x^\beta \frac{Z^\beta}{c^\beta} \leq |\Theta_x| Z_x^\beta \max_{\theta \in \Theta_x} e^{-\beta c(\theta)} w^0(\theta)$$

and thus using the bound in point b.

$$\min_y p_{x,y}^0 \min_{\theta \in \Theta_x} w^0(\theta) e^{-\beta \min_{\theta \in \Theta_x} c(\theta)} \leq m_x^\beta \frac{Z^\beta}{c^\beta} \leq |\Theta_x| \max_{\theta \in \Theta_x} w^0(\theta) e^{-\beta \min_{\theta \in \Theta_x} c(\theta)}$$

This implies that  $m_x^\beta$  concentrates (exponentially fast as  $\beta \rightarrow \infty$ ) on the point  $x$  that carries the **minimal spanning arborescence** for the cost  $c(x, y)$ , that is on the root of the tree  $\theta$  that minimizes  $c(\theta)$ . Indeed, for such an  $x$  and  $y \neq x$ ,  $m_y^\beta / m_x^\beta \rightarrow 0$  as  $\beta \rightarrow \infty$ .

d. We want to show that in the case  $c(x, y) = (V(y) - V(x))^+$  the root of the minimal spanning arborescence is the minimizer of  $V(\cdot)$  (which is unique since we assumed  $V$  injective). One may give a combinatorial proof in this case using  $c(x, y) - c(y, x) = V(y) - V(x)$ . Instead, let us give a *Markov Chain* proof, making good use of the fact that the minimal spanning arborescence does not depend on the values of  $p_{x,y}^0$ , but just on the edges on which this probability is non-vanishing.

First assume that the invariant measure  $m^0$  of  $(p_{x,y}^0)$  is reversible, then  $m_x^\beta := c^\beta e^{-\beta V(x)} Z_x^\beta m_x^0$  is reversible for  $(p_{x,y}^\beta)$  since

$$m_x^\beta p_{x,y}^\beta = Z_x^\beta e^{-\beta V(x)} \frac{e^{-\beta(V(y)-V(x))^+}}{Z_x^\beta} m_x^0 p_{x,y}^0 = e^{-\beta \max(V(x), V(y))} m_x^0 p_{x,y}^0 = e^{-\beta \max(V(x), V(y))} m_y^0 p_{y,x}^0 = m_y^\beta p_{y,x}^\beta$$

Thus, since  $V(x) + (V(y) - V(x))^+ = \max(V(y), V(x))$ , reasoning as above

$$m_x^\beta = c^\beta m_x^0 \sum_y e^{-\beta \max(V(y), V(x))} p_{x,y}^0$$

concentrates on the points  $x$  minimizing

$$x \mapsto \min_{y: p_{x,y}^0 > 0} \max(V(y), V(x))$$

In particular, if  $p_{x,x}^0 > 0$ , it concentrates on the minimizer of  $V(x)$ .

If  $m^0$  is not reversible, we can change  $p_{x,y}^0$  to  $q_{x,y}^0 := p_{x,y}^0 + p_{y,x}^0 m_y^0 / m_x^0$  (for which  $m^0$  is reversible) without changing the minimal spanning arborescence (which only depends on the  $c(x, y)$ ), so the answer does not change even when  $m^0$  is not reversible.

**Exercise 0.3.** Consider an irreducible Markov Chain on a finite state space  $S$  with transition probabilities  $(p_{x,y})$ . Recall that we consider  $S$  as a graph where the edges are the  $e = (x, y)$  such that  $p_{x,y} > 0$ . The weight of a **rooted spanning in-forest**  $F$  is the product

$$w(F) := \prod_{e \in F} p_e$$

For  $A \subset S$  nonempty, let  $\mathfrak{F}_A$  be the set of rooted spanning in-forests, whose set of roots is exactly  $A$ . For  $x \in S$  and  $r \in A$ , define  $\mathfrak{F}_A(x \rightarrow r)$  as the set of forests  $F \in \mathfrak{F}_A$  such that  $x$  is in an in-arborescence rooted at  $r$  (in other words, following the unique path in  $F$  outgoing from  $x$ , one reaches  $r \in A$ ).

Let now  $r \in A \subset S$ . Since the chain is irreducible and finite,  $\mathbb{P}(\tau_A < \infty) = 1$ , and let  $h(x) \equiv h_{r,A}(x) := \mathbb{P}_x(X_{\tau_A} = r)$ . Prove that

$$h(x) = \frac{\sum_{F \in \mathfrak{F}_A(x \rightarrow r)} w(F)}{\sum_{F \in \mathfrak{F}_A} w(F)} \quad (1)$$

### Solution

Since the Markov Chain is irreducible and has an invariant measure (being the state finite),  $h$  is the **unique** solution to the problem for the unknown  $u$

$$(I - \mathbf{1}_{A^c} P)u = \mathbf{1}_r$$

Let  $g(x)$  be the right hand side of Equation 1, we show that  $g$  also shows this equation, so  $g = h$ .

Clearly  $g(x) = \mathbf{1}_r(x)$  for  $x \in A$ , so we only need to check that  $g(x) = (Pg)(x)$  for  $x \notin A$ . The denominator  $\sum_{F \in \mathfrak{F}_A} w(F)$  simplifies on both sides of this equation so it remains to verify

$$\sum_{G \in \mathfrak{F}_A(x \rightarrow r)} w(G) \stackrel{?}{=} \sum_y p_{x,y} \sum_{F \in \mathfrak{F}_A(y \rightarrow r)} w(F), \quad x \notin A$$

Since  $\sum_y p_{x,y} = 1$  can we write this as

$$\sum_{z \neq x} \sum_{G \in \mathfrak{F}_A(x \rightarrow r)} p_{x,z} w(G) \stackrel{?}{=} \sum_{y \neq x} \sum_{F \in \mathfrak{F}_A(y \rightarrow r)} p_{x,y} w(F), \quad x \notin A$$

We can match terms on both sides, building a bijection between indices  $(z, G)$  (with  $x \rightarrow r$  in  $G$ ) on the left and  $(y, F)$  (with  $y \rightarrow r$  in  $F$ ) on the right as follows.

- If  $z \rightarrow r$  in  $G$ , then we take  $y = z$  and  $G = F$ , in other words these terms are on both sides.
- If  $z \nrightarrow r$ , then  $(x, z)$  is not in  $G$ . We match such a  $(z, G)$  with the element  $(y, F)$  in the r.h.s., where  $y$  is the unique vertex such that  $(x, y) \in G$ ; and  $F$  is the forest obtained from  $G$  removing  $(x, y)$  and adding  $(x, z)$ .

This is a bijection between summands in the l.h.s. and r.h.s. and clearly  $p_{x,z}w(G) = p_{x,y}w(F)$ .

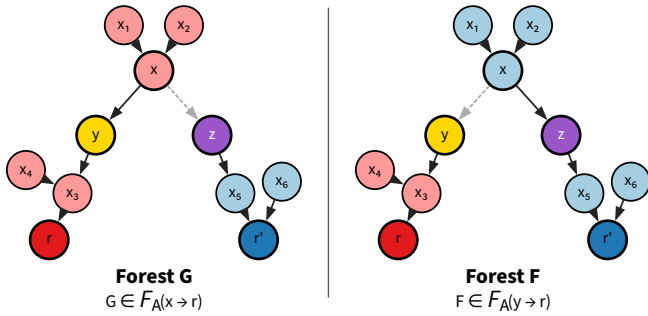


Figure 1: A pair  $(z, G)$  with  $G \in \mathfrak{F}_A(x \rightarrow r)$  and  $z \nrightarrow r$ , is mapped bijectively to a  $(y, F)$  with  $F \in \mathfrak{F}_A(y \rightarrow r)$ , just switching the successor of  $x$ .